Part IV

Database Design
Database Design

1. Phases of Database Design
2. Further Steps During Design
3. Capacity-preserving Transformations
4. ER-to-RM Transformation
Educational Objective for Today . . .

- Goals and steps of the database design process
- Rules to transform ER schemata into relational schemata
Goal of Design

- Data management for multiple application systems, for multiple years
- Therefore: special importance
- Design requirements
  - For every application, it should be possible to derive application data from data in the database — efficiently
  - Only store “sensible” (actually needed) data
  - Avoid redundancies
Phase Model

1. Requirement Analysis
2. Conceptional Design
3. Distribution Design
4. Logical Design
5. Data Definition
6. Physical Design
7. Implementation & Maintenance
Requirements Analysis

**Approach:** Collecting information needs from all specialist divisions

**Result:**
- Informal description (text, tabular lists, forms, etc.) of the problem domain
- Separation of the information about data (data analysis) from the information about functions (functional analysis)

"Classical" DB design:
- Only data analysis and following steps

Functional design:
- See methods of software engineering
Conceptual Design

- First formal description of the problem domain
- **Language means:** semantical data model
- **Process:**
  - Modeling of views, e.g., for different specialist divisions
  - Analysis of existing views with respect to conflicts
  - Integration of views into a full schema
- **Result:** full conceptual schema, e.g., ER diagram
Phases of Conceptual Design

- View Design
- View Analysis
- View Integration
Further Steps During Design

- ER modeling of different views of the complete information, e.g., for different specialist divisions of a company → conceptual design
  - Analysis and integration of views
  - Result: full conceptual schema
- Distribution design when using distributed storage
- Transformation to concrete implementation model (e.g., relational model) → logical design
- Data definition, implementation and maintenance → physical design
View Integration

- Analysis of existing view with respect to conflicts
- Integration of views into a full schema
Integration Conflicts

- **Naming conflicts**: Homonyms / synonyms
  - Homonyms: bank (money / river); order (command / request for goods)
  - Synonyms: car, vehicle, automobile

- **Typing conflicts**: different structures for the same element

- **Domain mismatch**: different domains for an element

- **Identifier conflicts**: e.g., different keys for the same element

- **Structural conflicts**: same fact expressed in different ways
Distribution Design

- If data should be distributed to several machines, a way of distributed storage must be determined.
- E.g., for a relation

  CUSTOMER (CNo, Name, Address, Zipcode, Account)

  - **Horizontal** distribution:
    CUSTOMER_1 (CNo, Name, Address, Zipcode, Account)
    \[ \text{where } \text{Zipcode} < 50\,000 \]
    CUSTOMER_2 (CNo, Name, Address, Zipcode, Account)
    \[ \text{where } \text{Zipcode} \geq 50\,000 \]

  - **Vertical** distribution (connection via attribute CNo):
    CUSTOMER_Adr (CNo, Name, Address, Zipcode)
    CUSTOMER_Account (CNo, Account)
Logical Design

- **Language means:** Data model of the chosen “implementation” DBMS, e.g., relational model

- **Process:**
  1. (Automatical) transformation of the conceptual schema, e.g., ER → relational model
  2. Improvement of the relational schema based on quality criteria (normalization, see Chapter 6):
     Design goals: avoid redundancies, ...

- **Result:** logical schema, e.g., collection of relation schemata
Data Definition

- Translation of logical schema into a concrete schema
- **Language means:** DDL and DML of DBMS (e.g., Oracle, DB2, SQL Server)
  - Database declaration in the DDL of the DBMS
  - Realization of integrity constraints
  - Definition of views
Physical Design

- Supplement physical design with support for efficient access, e.g., by defining indexes

Index
  - Access path: data structure for additional, key-based access to tuples (⟨key attribute value, tuple address⟩)
  - Usually implemented as a B*-tree

Language means: storage structure (definition) language SSL
Indexes in SQL

```sql
create [ unique ] index indexname
    on relname ( 
        attrname [ asc | desc ],
        attrname [ asc | desc ],
        ...
    )

Example

create index WineIdx on WINES (Name)
```
Necessity of Access Paths

- Example: Table with 100 GB of data, hard disk transfer rate of ca. 50 MB/s
- Operation: Search for a tuple (selection)
- Implementation: Sequential search
- Cost: $102.400/50 = 2.048$ sec. $\approx 34$ min.
Implementation and Maintenance

Phases of . . .

- Maintenance,
- Further optimization of the physical layer,
- Adaptation to new requirements or operating system platforms,
- Porting to new database management systems
- etc.
Transformation of the Conceptual Schema

- Translation to logical schema
  - Example: ER → RM
  - Correct?
  - Quality of transformation?

- Preservation of information capacity
  - Is it possible, after the transformation, to store exactly the same data as before?
  - . . . or more?
  - . . . or less?
Capacity-increasing Transformation

Transformation into

\[ R = \{\text{LicenseNo, Vineyard}\} \]

with exactly one key

\[ K = \{\{\text{LicenseNo}\}\} \]

Possible invalid relation:

<table>
<thead>
<tr>
<th>Has</th>
<th>LicenseNo</th>
<th>Vineyard</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>007</td>
<td>Helena</td>
</tr>
<tr>
<td></td>
<td>42</td>
<td>Helena</td>
</tr>
</tbody>
</table>
Capacity-preserving Transformation

- Correct instantiation

<table>
<thead>
<tr>
<th>Has</th>
<th>LicenseNo</th>
<th>Vineyard</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>007</td>
<td>Helena</td>
</tr>
<tr>
<td></td>
<td>42</td>
<td>Müller</td>
</tr>
</tbody>
</table>

- Correct set of keys

\[ K = \{\{\text{LicenseNo}\}, \{\text{Vineyard}\}\} \]
Capacity-decreasing Transformation

- Relation schema with one key \{WName\}
- Instantiation that is no longer valid:

<table>
<thead>
<tr>
<th>ConsistsOf</th>
<th>WName</th>
<th>GrapeName</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Zinfandel Red Blossom</td>
<td>Zinfandel</td>
</tr>
<tr>
<td></td>
<td>Bordeaux Blanc</td>
<td>Cabernet Sauvignon</td>
</tr>
<tr>
<td></td>
<td>Bordeaux Blanc</td>
<td>Muscadelle</td>
</tr>
</tbody>
</table>

- Capacity-preserving when using the keys of both entity types as the new key in the relation schema

\[ K = \{\{WName, GrapeName\}\} \]
Example Transformation ER-RM: Input

Database Design
ER-to-RM Transformation

Example Transformation ER-RM: Input

Database Concepts
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4–22
Example Transformation ER-RM: Result

1. \( \text{GRAPE} = \{\text{Color}, \text{GrapeName}\} \) with \( K_{\text{GRAPE}} = \{\{\text{GrapeName}\}\} \)
2. \( \text{ConsistsOf} = \{\text{GrapeName}, \text{WName}, \text{Amount}\} \) with \( K_{\text{ConsistsOf}} = \{\{\text{GrapeName}, \text{WName}\}\} \)
3. \( \text{WINE} = \{\text{Color}, \text{WName}, \text{Vintage}, \text{Res.Sugar}\} \) with \( K_{\text{WINE}} = \{\{\text{WName}\}\} \)
4. \( \text{PRODUCE} = \{\text{WName}, \text{Vineyard}\} \) with \( K_{\text{PRODUCE}} = \{\{\text{WName}\}\} \)
5. \( \text{PRODUCER} = \{\text{Vineyard}, \text{Address}\} \) with \( K_{\text{PRODUCER}} = \{\{\text{Vineyard}\}\} \)
ER Transformation into Relations

- **Entity types and relationship types**: both transformed into relation schemata
- **Attributes**: attributes of the relation schema, **keys** are adopted
- **Cardinalities** of the relationships: expressed in respective relation schemata by choice of keys
- In some cases: **merge** of the relation schemata of entity and relationship types
- Introduce foreign key constraints between the remaining relation schemata
Transformation of Relationship Types

- New relation schema with all attributes of the relationship type; additionally, adopt all primary keys of the participating entity types

- **Determining keys:**
  - **m:n relationship:** both primary keys together form the key in the new relation schema
  - **1:n relationship:** primary keys of the n-side (in the functional notation, this is the side without the arrowhead) form key in the new relation schema
  - **1:1 relationship:** both primary keys become a key in the new relation schema; the primary key is then chosen from these keys
n:m Relationships

Transformation

1. GRAPE = \{\text{Color, GrapeName}\} with \(K_{\text{GRAPE}} = \{\{\text{GrapeName}\}\}\)
2. ConsistsOf = \{\text{GrapeName, WName, Amount}\} with \(K_{\text{ConsistsOf}} = \{\{\text{GrapeName, WName}\}\}\)
3. WINE = \{\text{Color, WName, Vintage, Res. Sugar}\} with \(K_{\text{WINE}} = \{\{\text{WName}\}\}\)

Attributes \text{GrapeName} and \text{WName} together are key
1:n Relationships

(Preliminary) transformation

- PRODUCER with the attributes Vineyard and Address,
- AREA with the attributes Name and Region, and
- LocatedIn with the attributes Vineyard and Name and the primary key of the $n$-side Vineyard as primary key of this schema.
Possible Merges

- **Optional relationships** ([0,1] or [0,n]) are not merged
- With cardinalities [1,1] or [1,n] (**mandatory relationships**), merge is possible:
  - 1:n relationship: the entity-relation schema of the n-side can be integrated into the relation schema of the relationship
  - 1:1 relationship: both entity-relation schemata can be integrated into the relation schema of the relationship
1:1 Relationships

(Preliminary) transformation

- PRODUCER with the attributes Vineyard and Address
- LICENSE with the two attributes LicenseNo and Hectoliters
- Has with the primary keys of both participating entity types each as key of this schema, that is LicenseNo and Vineyard
### 1:1 Relationships: Merge

**Transformation with merge**
- **Merged relation:**

<table>
<thead>
<tr>
<th>PRODUCER</th>
<th>Vineyard</th>
<th>Address</th>
<th>LicenseNo</th>
<th>Hectoliters</th>
</tr>
</thead>
<tbody>
<tr>
<td>Rotkäppchen</td>
<td>Freiberg</td>
<td>42-007</td>
<td>10 000</td>
<td></td>
</tr>
<tr>
<td>Vineyard Müller</td>
<td>Dagstuhl</td>
<td>42-009</td>
<td>250</td>
<td></td>
</tr>
</tbody>
</table>

- **Producers without license require null values:**

<table>
<thead>
<tr>
<th>PRODUCER</th>
<th>Vineyard</th>
<th>Address</th>
<th>LicenseNo</th>
<th>Hectoliters</th>
</tr>
</thead>
<tbody>
<tr>
<td>Rotkäppchen</td>
<td>Freiberg</td>
<td>42-007</td>
<td>10 000</td>
<td></td>
</tr>
<tr>
<td>Vineyard Müller</td>
<td>Dagstuhl</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

- **Free Licenses lead to additional null values:**

<table>
<thead>
<tr>
<th>PRODUCER</th>
<th>Vineyard</th>
<th>Address</th>
<th>LicenseNo</th>
<th>Hectoliters</th>
</tr>
</thead>
<tbody>
<tr>
<td>Rotkäppchen</td>
<td>Freiberg</td>
<td>42-007</td>
<td>10 000</td>
<td></td>
</tr>
<tr>
<td>Vineyard Müller</td>
<td>Dagstuhl</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>42-003</td>
<td>100 000</td>
</tr>
</tbody>
</table>
Dependent Entity Types

Transformation

1. \( \text{WINEVINTAGE} = \{\text{WName, Res. Sugar, Year}\} \) with \( K_{\text{WINEVINTAGE}} = \{\{\text{WName, Year}\}\} \)

2. \( \text{WINE} = \{\text{Color, WName}\} \) with \( K_{\text{WINE}} = \{\{\text{WName}\}\} \)

- Attribute \text{WName} in \text{WINEVINTAGE} is foreign key to relation \text{WINE}
IS-A Relationship

Transformation

1. \( \text{WINE} = \{\text{Color, WName}\} \) with \( K_{\text{WINE}} = \{\{\text{WName}\}\} \)
2. \( \text{SPARKLING\_WINE} = \{\text{WName, Production}\} \) with \( K_{\text{SPARKLING\_WINE}} = \{\{\text{WName}\}\} \)

- \text{WName} in \( \text{SPARKLING\_WINE} \) is foreign key with respect to relation \( \text{WINE} \)
Recursive Relationships

Transformation

1. \( \text{AREA} = \{\text{Name, Region}\} \) with \( K_{\text{AREA}} = \{\{\text{Name}\}\} \)
2. \( \text{ADJOINS} = \{\text{to, from}\} \) with \( K_{\text{ADJOINS}} = \{\{\text{to, from}\}\} \)
Recursive Functional Relationships

**Transformation**

1. \( \text{CRITIC} = \{\text{Name, Organization, Mentorname}\} \) with

   \( K_{\text{CRITIC}} = \{\{\text{Name}\}\} \)

   - Mentorname is foreign key to attribute Name of relation CRITIC.
N-ary Relationships

- Every participating entity type is treated according to the rules stated above.
- For relationship `Recommends`, the primary keys of the three participating entity types are included in the resulting relation schema.
- Relationship has a generic type \((k:m:n)\) relationship: all primary keys together form the key.
N-ary Relationships: Result

1. \( \text{RECOMMENDS} = \{ \text{WName}, \text{DName}, \text{Name} \} \) with \( K_{\text{RECOMMENDS}} = \{ \{ \text{WName}, \text{DName}, \text{Name} \} \} \)
2. \( \text{DISH} = \{ \text{DName}, \text{Side\textunderscore Dish} \} \) with \( K_{\text{DISH}} = \{ \{ \text{DName} \} \} \)
3. \( \text{WINE} = \{ \text{Color}, \text{WName}, \text{Vintage}, \text{Res\textunderscore Sugar} \} \) with \( K_{\text{WINE}} = \{ \{ \text{WName} \} \} \)
4. \( \text{CRITIC} = \{ \text{Name}, \text{Organization} \} \) with \( K_{\text{CRITIC}} = \{ \{ \text{Name} \} \} \)

- The three key attributes of \( \text{RECOMMENDS} \) are foreign keys to the respective source relations (\( \text{CRITIC, WINE, DISH} \)).
# Overview of Transformations

<table>
<thead>
<tr>
<th>ER Concept</th>
<th>Is Translated into Relational Concept</th>
</tr>
</thead>
<tbody>
<tr>
<td>Entity type $E_i$</td>
<td>Relation schema $R_i$</td>
</tr>
<tr>
<td>Attributes of $E_i$</td>
<td>Attributes of $R_i$</td>
</tr>
<tr>
<td>Primary key $P_i$</td>
<td>Primary key $P_i$</td>
</tr>
<tr>
<td>Relationship type</td>
<td>Relation schema</td>
</tr>
<tr>
<td>Its attributes</td>
<td>Attributes: $P_1$, $P_2$</td>
</tr>
<tr>
<td>$1 : n$</td>
<td>Further attributes</td>
</tr>
<tr>
<td>$1 : 1$</td>
<td>$P_2$ becomes primary key of the relationship</td>
</tr>
<tr>
<td>$m : n$</td>
<td>$P_1$ and $P_2$ become key of the relationship</td>
</tr>
<tr>
<td>IS-A relationship</td>
<td>$P_1 \cup P_2$ becomes primary key of the relationship</td>
</tr>
<tr>
<td></td>
<td>$R_1$ gets an additional key $P_2$</td>
</tr>
</tbody>
</table>

$E_1$, $E_2$: Entity types participating in a relationship,
$P_1$, $P_2$: Their primary keys,
$1 : n$ relationship: $E_2$ is $n$-side,
IS-A relationship: $E_1$ is a special entity type
Summary

- Phases of database design
- Capacity-preserving transformations
- Transformation ER $\rightarrow$ relational
Control Questions

- Which steps does the database design process comprise?
- Which requirements do the transformations between each design step have to fulfill? Why?
- How are concepts of the ER model translated into concepts of the relational model?
- How are the different cardinalities of relationship types accounted for during transformation?